

Goal: Motivate & explain some constructions.

(Those who were in Bonn might recognise some ideas from William's talk: I would say William's ideas are generalizations of these)

§1. General Adams sseq's

Def. \mathcal{C} stable ∞ -cat, \mathcal{A} ab. cat.
A **homological functor** is $h: \mathcal{C} \rightarrow \mathcal{A}$ s.t.

$$X \rightarrow Y \rightarrow Z \text{ cofib.} \Rightarrow h(X) \rightarrow h(Y) \rightarrow h(Z)$$

Let \mathcal{A} be equipped with an auto-equiv. $[i]: \mathcal{A} \rightarrow \mathcal{A}$.
A **homology theory** $\mathcal{C} \rightarrow (\mathcal{A}, [i])$ is homological functor h together with "suspension iso's".

$$h(\Sigma X) \cong h(X)[i].$$

Eg. $\mathcal{A} = \text{gr Ab}$, $[i]$ is grading shift: $(M[i])_n = M_{n-1}$.

Then $E \in \mathcal{S}_p$ yields hly thg $E_*(-): \mathcal{S}_p \rightarrow \text{gr Ab}$.

Def. A homological functor $h: \mathcal{C} \rightarrow \mathcal{A}$ is **adapted** if

1) \mathcal{A} has enough injectives.

2) $\forall I \in \mathcal{A}$ injective, $\exists X_I \in \mathcal{C}$ s.t.

$$[-, X_I]_{\mathcal{C}} \cong \text{Hom}_{\mathcal{A}}(h(-), I).$$

if exists, then determined uniquely up to isom. by Yoneda

3) $h(X_I) \cong I$.

Remk. (1) & (2) is usually formal (Brown representability).
(3) is very not formal: that's where the context is.

Idea. h adapted means: \mathcal{C} has an h -ASS where E_2 page is computed by $\text{Ext}_{\mathcal{A}}$.

Eg. $A \in \text{CAlg}(\mathcal{S}_p)$. Then $\pi_*: \text{Mod } A \rightarrow \text{Mod } A_*$ is adapted.

If I injective, then by Brown $\exists X_I$ as in (2). But:

$$\pi_n X_I = [\Sigma^n A, X_I]_A \cong \text{Hom}_{A_*} (A_*[n], I) \cong I_n.$$

So $\pi_* X_I \cong I$. Resulting ASS is the **Ext-sseq**.

⚠ Adaptedness can be destroyed by precomposing with other functors.

Eg $H_*(-; \mathbb{F}_p)$ is $S_p \rightarrow \text{Mod}_{\mathbb{F}_p}(S_p) \xrightarrow{\pi_*} \text{grVect}_{\mathbb{F}_p}$.

Then $\mathbb{F}_p \in \text{grVect}_{\mathbb{F}_p}$ is injective, and \mathbb{F}_p is injective left:

$$\text{Hom}_{\mathbb{F}_p}(H_*(X; \mathbb{F}_p), \mathbb{F}_p) \cong H^*(X; \mathbb{F}_p) \cong [X, \mathbb{F}_p].$$

But: $H_*(\mathbb{F}_p; \mathbb{F}_p) = \mathcal{A}_*$ is very different from $\pi_*(\mathbb{F}_p) = \mathbb{F}_p$.

Reason this fails: $H_*(X; \mathbb{F}_p)$ as \mathbb{F}_p -vs. does not remember enough about X . Need to "add information" $\leadsto \mathcal{A}_*$ -comodules.

$$\begin{array}{ccc} S_p & \longrightarrow & \text{Mod}_{\mathbb{F}_p}(S_p) \\ H_*(-; \mathbb{F}_p) \downarrow & & \downarrow \pi_* \\ \text{Comod}_{\mathcal{A}_*} & \longrightarrow & \text{grVect}_{\mathbb{F}_p} \end{array}$$

Thm: E Adams type. Then $E_*(-): S_p \rightarrow \text{Comod}_{E_*E}$ is adapted.

This leads to the E -ASS as before. (This is not nec. the one from $E^{(0,1)}$)

Thm: Let $h: \mathcal{C} \rightarrow \mathcal{A}$ be homological functor. Then \exists adapted $H: \mathcal{C} \rightarrow \mathcal{B}$ together with an exact functor $\mathcal{B} \rightarrow \mathcal{A}$ and

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{H} & \mathcal{B} \\ & \searrow h & \downarrow \\ & & \mathcal{A} \end{array}$$

Moreover, \mathcal{B} is comonadic over \mathcal{A} for a left exact comonad.

Eg. $\text{Comod}_{\mathcal{A}_*} \cong \text{Comod}_{\mathcal{C}}(\text{grVect}_{\mathbb{F}_p})$ for $\mathcal{C}: V \mapsto \mathcal{A}_* \otimes_{\mathbb{F}_p} V$.

Rmk: • This (\mathcal{B}, H) is unique if it exists: adapted hghy th's form a poset
• The proof does not help in finding an explicit \mathcal{B} in practise.

Later: can build a "synthetic category" $\text{Syn}_h(\mathcal{C})$ when h adapted.

§2. Synthetic A -modules (Take it slow: split it up into two steps.)

$A \in \text{Cat}_{\mathbb{A}}$. Define synthetic cat for the Ext seq for Mod_A :

Def: $\text{Syn}_A(\text{Mod}_A) := \text{PSh}_{\Sigma}(\text{Mod}_A^{\text{ff}}; S_p)$.
presheaves preserving finite products \leadsto finite & free \leftarrow closed under \otimes in $\text{Mod}_A \Rightarrow$ get Day convolution

Think: $\text{Syn}_A(\text{Mod}_A)$ is associated to adapted $\pi_*: \text{Mod}_A \rightarrow \text{Mod}_{A_*}(\text{grAb})$.

Gets t-structure: use levelwise t-str. using standard one on Sp .

Then

$$\begin{aligned} \text{Syn}_A(\text{Mod}_A) &\simeq \text{PSh}_\Sigma(\text{Mod}_A^{\text{ff}}; \text{Set}) \\ &\simeq \text{PSh}_\Sigma(\text{Mod}_A^{\text{ff}}; \text{Ab}) \\ &\simeq \text{PSh}_\Sigma(\text{hMod}_A^{\text{ff}}; \text{Ab}) \\ &\simeq \text{Mod}_{A_*}(\text{grAb}) \end{aligned}$$

because we're looking at prod.-pres. presheaves.

(this sounds fancy but it's not)

$\text{hMod}_A \simeq \text{Mod}_{A_*}(\text{grAb})$ is even an equivalence of lax theories

Def: $v: \text{Mod}_A \rightarrow \text{Syn}_A(\text{Mod}_A)$ is $M \mapsto \text{Map}(-, M)$.

Here we use: $\text{PSh}_\Sigma(\text{Mod}_A^{\text{ff}}; \mathcal{S}) \simeq \text{PSh}_\Sigma(\text{Mod}_A^{\text{ff}}; \text{Sp}_{\geq 0}) \subseteq \text{PSh}_\Sigma(\text{Mod}_A^{\text{ff}}; \text{Sp})$.

Then v lands in $\text{Syn}^{\geq 0}$. Hence:

$$\text{Mod}_A \xrightarrow{v} \text{Syn}^{\geq 0} \xrightarrow{\tau_{\leq 0}} \text{Syn}^\heartsuit \simeq \text{Mod}_{A_*}(\text{grAb}) \text{ is } \pi_*$$

Could say: I'm not gonna explain everything in terms of this construction. I already showed how everything can be done in terms of the abstract properties.

But: $\text{PSh}_\Sigma(\text{Mod}_{A_*}^{\text{ff}}; \text{Sp}) \simeq \mathcal{D}(A)$. ∇

Prop: v is fully faithful.

Proof: fully faithful on Mod_A^{ff} by Tonda. But also preserves filtered colimits. \square

§3. Synthetic Spectra

We essentially have to modify the previous setting to "make it adapted" again.

$$\begin{array}{ccc} \text{Sp} & \xrightarrow{E_0^-} & \text{Mod}_E \\ \downarrow & & \downarrow \pi_* \\ \text{Comod}_{E_*E} & \rightarrow & \text{Mod}_{E_*} \end{array}$$

← this picture is only true if E is highly structured.

Let E be htyg assoc. ring sp.

Def: $\text{Sp}_E^{\text{fp}} \subseteq \text{Sp}^{\omega}$ those P s.t. E_*P projective (\Leftrightarrow dualisable).
 $\text{Equip}_{\text{Sp}_E^{\text{fp}}}$ with a designated class of "epimorphisms": the E_* -surjections.

Def: $\text{Syn}_E \subseteq \text{PSh}_\Sigma(\text{Sp}_E^{\text{fp}}; \text{Sp})$ full subcat. on those F s.t.

$$\textcircled{*} \quad X \rightarrow Y \rightarrow Z \text{ cofib seq. of spectra in } \text{Sp}_E^{\text{fp}} \quad \text{with } E_*Y \rightarrow E_*Z \xrightarrow{\text{surj}} \Rightarrow F(Z) \rightarrow F(Y) \rightarrow F(X) \text{ cofib seq.}$$

It's very convenient to be able to phrase this in the language of sheaves.

Make Sp_E^{fp} a site. a cover of T is $\{Q \rightarrow T\}$ single map with $E_* Q \rightarrow E_* T$ surj.

sheaf condition is similar to comonad in the adapted factorisation

Thm (Pstragowski): $\mathcal{O} \Leftrightarrow F$ a sheaf for this topology.

Remark we really have a functor $Alg(hSp) \rightarrow \mathcal{E}sites$, so also get a functor $Syn_{(-)}(Sp) : Alg(hSp) \rightarrow Prst$. (in fact, have both adjoints)

Thm (Goerss-Hopkins): (A, T) Hopf algebra of Adams type. Then $Comod_{(A, T)} \cong Sh_{\Sigma}(Comod_{(A, T)}^{fp}; Set)$. coverings are surjections

Thm (Pstragowski): Let E be Adams type. Then

Also: $Sp \xrightarrow{\sim} Syn_E^{fp} \xrightarrow{\sim} Syn_E^{\vee}$ is $E_*(-)$.

$$Sh_{\Sigma}(Comod_{E_*E}^{fp}; Set) \cong Sh_{\Sigma}(Sp_E^{fp}; Set)$$

induced by $E_*(-) : Sp_E^{fp} \rightarrow Comod_{E_*E}^{fp} \Rightarrow$ only for those E do we understand Syn_E as defined like this

But: sheaves of spectra are different:

$$Sh_{\Sigma}(Comod_{E_*E}^{fp}; Sp) \cong Stable_{E_*E} \xleftarrow{\sim} Mod_{CC}(Syn_E Sp)$$

"Algebra can only reconstruct the E_2 -page".

(but this inclusion is subtle)

Still have $Sh_{\Sigma}(-; S) \cong Sh_{\Sigma}(-; Sp_{20}) \subseteq Sh_{\Sigma}(-; Sp)$, so can define

Def: $v : Sp \rightarrow Syn_E(Sp)$, $X \mapsto Map(-, X)$.

This is fully faithful by the same argument.

Turns out: hypercomplete $\Leftrightarrow vE$ -local (hypercomplete: receiving no maps from ω -conn.)

§4. Filtered models (Essentially in (BHS, App C), but this slightly different perspective will be in (Barkan-UN.))

Recall: $FilSp \xrightleftharpoons[\mathcal{F}]{\mathcal{P}} Syn_E$.

Thm: This is monadic $\Leftrightarrow Syn_E$ is generated under colimits by $\{\mathcal{S}^{n,s}\}_{n,s}$. If this happens, we say Syn_E is cellular. \rightsquigarrow [Lawson], [Barkan-UN].

Remains to identify the monad on $FilSp$. Can identify $\sigma(vS)$, but we need to identify this as an E_* -alg. For that, need lots of structure on E :

Prop: If $E \triangleright E_{\omega}$ and E -ASS(S) conv. cond., then $Syn_E \cong Mod_{ASS_E(S)}(FilSp)$. Eg. MU.

§4. Recollection on derived ω -cats.

Def: A **prestable** ω -cat. is an ω -cat. \mathcal{C} s.t. \exists stable \mathcal{D} with t -str. s.t. $\mathcal{C} \cong \mathcal{D}_{\geq 0}$.

A **Grothendieck prestable** ω -cat. is a prestable one that's presentable and where filtered colimits are left exact.

If \mathcal{C} is prestable and $\mathcal{C} = \mathcal{D}_{\geq 0}$, then notion of "n-truncated obj." in \mathcal{C} agrees with the t -structure one on \mathcal{D} . $\mathcal{C}_{\leq n} \cong (\mathcal{D}_{\geq 0})_{\leq n}$.
 Conic $\mathcal{C} := \mathcal{C}_{\leq 0}$. ("Right complete t -str" $\Leftrightarrow \mathcal{D} = \text{Sp}(\mathcal{D}_{\geq 0}) = \text{Sp}(\mathcal{C})$)

Proto-theorem: if \mathcal{A} ab. cat., \mathcal{E} prestable, then

$$\text{Fun}^{\text{ex}}(\mathcal{D}_{\geq 0}(\mathcal{A}), \mathcal{E}) \cong \text{Fun}^{\text{ex}}(\mathcal{A}, \mathcal{E}^{\heartsuit}).$$

To make precise, need some hypotheses:

$$\text{Fun}^{\text{ex}}(\mathcal{D}_{\geq 0}^b(\mathcal{A}), \mathcal{E}) \cong \text{Fun}^{\text{ex}}(\mathcal{A}, \mathcal{E}^{\heartsuit}). \quad \mathcal{E} \text{ finite limits [FP, Thm. 5.10]}$$

$$\text{LFun}^{\text{ex}}(\check{\mathcal{D}}_{\geq 0}(\mathcal{A}), \mathcal{E}) \cong \text{LFun}^{\text{ex}}(\mathcal{A}, \mathcal{E}^{\heartsuit}) \quad \mathcal{E} \text{ Grothendieck}$$

$$\text{LFun}^{\text{ex}}(\mathcal{D}_{\geq 0}(\mathcal{A}), \mathcal{E}) \cong \text{LFun}^{\text{ex}}(\mathcal{A}, \mathcal{E}^{\heartsuit}) \quad \mathcal{E} \text{ Grothendieck \& separated [SAG, Thm. C.5.4.9]}$$

$$\mathcal{D}(\mathcal{A}) = \check{\mathcal{D}}(\mathcal{A}) / \{ \omega\text{-connective objects} \}. \quad \text{"hypercompletion"}$$

Construction [Patchkora-Pstragowski]: For \mathcal{E} stable, compactly gen., $h: \mathcal{C} \rightarrow \mathcal{A}$ adapted, \mathcal{A} Grothendieck abelian, h pres. arbitrary sums

$\mapsto \check{\mathcal{D}}(\mathcal{C}; h)$ stable ω -cat. with right complete t -str.

$$\check{\nu}: \mathcal{C} \rightarrow \check{\mathcal{D}}_{\geq 0}(\mathcal{C}; h).$$

$$\mapsto \mathcal{D}(\mathcal{C}; h) := \check{\mathcal{D}}(\mathcal{C}; h) / \{ \omega\text{-com.} \}$$

Similar properties as $\text{Syn}_{\mathcal{E}}$, but notably:

1) no monoidal str. in general.

2) t -str. slightly less structured: in general, there's no object X s.t. (co)connectivity is measured by $\tau_{x, x}(X \otimes -)$.

maybe should emphasize this more in other notes

Still, $\check{\mathcal{D}}^{\heartsuit}(\mathcal{C}; h) \cong \mathcal{A}$, and even: $\mathcal{C} \xrightarrow{\check{\nu}} \mathcal{D}_{\geq 0}(\mathcal{C}; h) \xrightarrow{\tau_{\geq 0}} \mathcal{D}^{\heartsuit}(\mathcal{C}; h) \cong \mathcal{A}$
 \Downarrow
 $\hookrightarrow t\text{-str. http graph is homology with values in } \mathcal{A}.$

In fact, it's the universal such pair:

Thm: If E is Groth. prestable, then

$$G: \check{D}_{\neq 0}(E; h) \rightarrow E \quad \Leftrightarrow \quad G_0: \mathcal{A} \rightarrow E^\heartsuit + \begin{array}{ccc} \mathcal{G}_0 & \xrightarrow{h} & \mathcal{A} \xrightarrow{G_0} E^\heartsuit \\ & \nearrow & \downarrow \tau_{\leq 0} \\ & & E \end{array}$$

exact, colin-pres. exact, colin-pres. "prestable enhancement"

Construction: idea is: for Syn , the equiv. $Syn_\Sigma(Sp_E^{fp}; Set) \cong Comod_{E_*E}$ was key. To generalise this, would need to find a suitable version of this morph.

Universal case: $\mathcal{A} = PSh_\Sigma(E^\omega; Set)$ (think: ridiculous enlargement of π_*)

and $h: X \mapsto [-, X]_E$. (preserves \otimes by compactness)

Then define $\check{D}(E; h) := PSh_\Sigma(E^\omega; Sp)$

Compactly gen. case: $h: E \rightarrow \mathcal{A} \Leftrightarrow$ quotient $PSh_\Sigma(E^\omega; Ab) / \mathcal{K}$ localising subcat. \mathcal{K} .

Left \mathcal{K} to $PSh_\Sigma(E^\omega; Sp)$: take smallest loc. subcat. \mathcal{R} containing $\mathcal{K} \subseteq PSh_\Sigma(E^\omega; Sp)$.

Then define $D(E; h) := PSh_\Sigma(E^\omega; Sp) / \mathcal{R}$.

§ Synthetic spectra as a derived ω -cat.

Consider $E_*(-): Sp \rightarrow Comod_{E_*E} \rightsquigarrow \check{D}(Sp; E_*(-))$.

Prop: $Syn_E(Sp) \cong Ind\ Thick(\check{D}P \mid P \in Sp_E^{fp})$.

Δ $Syn \neq D$, but rather $Syn \cong D$ (but $Syn \neq \hat{D} \mathbb{P}S$)

Now, Sp_E^{fp} has monoidal str. Give $Thick(\check{D}P)$ a mon. str. by declaring \check{D} to be symmetric monoidal on $P \in Sp_E^{fp}$. Then Ind cat also gets a mon. str., and this is equiv. to (Syn_E, \otimes) .

Sometimes this becomes simpler: if E is a field, then simply

$$Syn_E \cong \check{D}(Sp; E_*(-)).$$

In part, such equivs give Syn_E a new univ. property.

Variant: version of E -ASS defined via $E^{\otimes \bullet+1}$. By categorical const., there is ab. cat. \mathcal{A} with $F: Sp \rightarrow \mathcal{A}$ adapted giving rise to this. But not clear how to describe this in general.

If $P \in Sp_E^{fp}$, then $Ext_{\mathcal{A}}(F(P), F(X)) \cong Ext_{E_*E}(E_*P, E_*X)$.