Hilbert 10 via additive combinatorics I

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Question (Hilbert's tenth problem)

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In modern terms: does there exist an algorithm such that: **Input:** a polynomial $p \in \mathbb{Z}[x_1, \ldots, x_n]$. **Output:** "YES" if there is an integer solution $(a_1, \ldots, a_n) \in \mathbb{Z}^n$ with $p(a_1, \ldots, a_n) = 0$, "NO" otherwise.

Definition (Diophantine set)

We say that a subset $S \subseteq \mathbb{Z}^n$ is Diophantine if there exists a polynomial $p(x_1, \ldots, x_n, y_1, \ldots, y_m) \in \mathbb{Z}[\mathbf{x}, \mathbf{y}]$ such that

 $S = \{ \mathbf{x} \in \mathbb{Z}^n : \exists \mathbf{y} \in \mathbb{Z}^m \text{ such that } p(\mathbf{x}, \mathbf{y}) = 0 \}.$

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Corollary (Hilbert's tenth problem)

Hilbert's tenth problem is undecidable, i.e. there is no algorithm that can decide whether a polynomial $p \in \mathbb{Z}[x_1, \ldots, x_n]$ has a zero or not.

In 1950, Julia Robinson proves that "J.R." implies that exponentiation is Diophantine, i.e.

$$\{(a,b,c)\in\mathbb{Z}^3:a=b^c\}$$

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Matiyasevich asks in the 1970s: what about other rings?

Definition

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This implies the previous theorem by work of Eisenträger.

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Theorem (Shlapentokh (2008))

Let L/K be an extension of number fields. Suppose that there exists an elliptic curve E/K such that $\operatorname{rk} E(L) = \operatorname{rk} E(K) > 0$. Then O_K is Diophantine over O_L .

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Our main new technical result is (rank growth theorem):

Theorem (K.-Pagano, 2024)

Let K be a number field with at least 32 real embeddings. Then there exists an elliptic curve E/K such that

 $\operatorname{rk} E(K) = \operatorname{rk} E(K(i)) > 0.$

Lemma (Denef, Denef–Lipshitz)

Let $K \subseteq L$ be number fields. Then: (i) If $D_1, D_2 \subseteq O_K$ are Diophantine over O_K , so is $D_1 \cap D_2$.

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Rank growth implies that \mathbb{Z} is Diophantine over O_K .

Let K be a number field: we will show that \mathbb{Z} is Dio over O_K . Define M to be the Galois closure of $K(i, \sqrt{3}, \sqrt{5}, \sqrt{7}, \sqrt{11}, \sqrt{13}, \sqrt{17})$, so suffices to show that \mathbb{Z} is Dio over O_M by (*iii*).

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Rank growth implies that \mathbb{Z} is Diophantine over $O_{\mathcal{K}}$.

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We recall:

Theorem (K.–Pagano, 2024)

Let K be a number field with at least 32 real embeddings. Then there exists an elliptic curve E/K such that

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- lt is also easy to use 2-descent to upper bound the rank of E/K.

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What is the challenge?

▶ It is easy to construct elliptic curves E/K with $\operatorname{rk} E(K) > 0$.

It is also easy to use 2-descent to upper bound the rank of E/K. However, it is not at all clear how to combine this.

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Taking Galois cohomology gives

$$0 \to E(K)[2] \to E(K) \xrightarrow{\cdot 2} E(K) \xrightarrow{\delta} H^1(G_K, E[2]),$$

so $E(K)/2E(K) \stackrel{\delta}{\hookrightarrow} H^1(G_K, E[2]).$

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Recall:

$$\begin{array}{ccc} E(K)/2E(K) & & \stackrel{\delta}{\longrightarrow} & H^1(G_K, E[2]) \\ & & & \downarrow^{\mathrm{res}_{\nu}} \\ E(K_{\nu})/2E(K_{\nu}) & & \stackrel{\delta_{\nu}}{\longrightarrow} & H^1(G_{K_{\nu}}, E[2]) \end{array}$$

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We define

$$\operatorname{Sel}^{2}(E/K) := \operatorname{ker}\left(H^{1}(G_{K}, E[2]) \xrightarrow{\prod_{v} \operatorname{res}_{v}} \prod_{v} \frac{H^{1}(G_{K_{v}}, E[2])}{\operatorname{im}(\delta_{v})}\right)$$

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Key facts:

(1) We have $im(\delta) \subseteq Sel^2(E/K)$, and hence

 $\operatorname{rk} E(K) + \dim_{\mathbb{F}_2} E(K)[2] = \dim_{\mathbb{F}_2} E(K)/2E(K) \leq \dim_{\mathbb{F}_2} \operatorname{Sel}^2(E/K).$

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 $\operatorname{rk} E(K) + \dim_{\mathbb{F}_2} E(K)[2] = \dim_{\mathbb{F}_2} E(K)/2E(K) \leq \dim_{\mathbb{F}_2} \operatorname{Sel}^2(E/K).$

(2) The group $\operatorname{Sel}^2(E/K)$ is computable and finite dimensional (note that rank is not known to be computable!).

Take some E_1 with $E_1(\mathcal{K})[2] \cong \mathbb{F}_2^2$, i.e.

$$E_1: y^2 = (x - a_1)(x - a_2)(x - a_3)$$
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Consider the quadratic twist E_1^t with $t := (n - a_1)(n - a_2)(n - a_3)$. Then

$$E_1^t$$
: $(n - a_1)(n - a_2)(n - a_3)y^2 = (x - a_1)(x - a_2)(x - a_3)$

has the rational point (x, y) = (n, 1).

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has the rational point (x, y) = (n, 1). It is not hard to show that this point is almost never torsion (Northcott's theorem + height bounds).

Take some E_1 with $E_1(K)[2] \cong \mathbb{F}_2^2$, i.e.

$$E_1: y^2 = (x - a_1)(x - a_2)(x - a_3)$$
 $a_1, a_2, a_3 \in K$ distinct.

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However, it is very difficult to control rank in this family. Instead, we consider E_1^t with $t := m(n - a_1m)(n - a_2m)(n - a_3m)$. Then

$$E_1^t: m(n-a_1m)(n-a_2m)(n-a_3m)y^2 = (x-a_1)(x-a_2)(x-a_3)$$

has the rational point $(x, y) = (n/m, 1/m^2)$ and hence $\operatorname{rk} E_1^t(K) > 0$.

Because $i \notin K$ (since K has 32 real places), we get

 $\operatorname{rk} E_1^t(K(i)) = \operatorname{rk} E_1^t(K) + \operatorname{rk} E_1^{-t}(K).$

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Hence it suffices to find $t := m(n - a_1m)(n - a_2m)(n - a_3m)$ with

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as then $E := E_1^t$ satisfies $\operatorname{rk} E(K(i)) = \operatorname{rk} E(K) > 0$ as desired.

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Because E_1 has full rational 2-torsion (and therefore all of its quadratic twists do), it suffices to find $t := m(n - a_1m)(n - a_2m)(n - a_3m)$ with

$$\dim_{\mathbb{F}_2} \mathrm{Sel}^2(E_1^{-t}/K) = 2,$$

as this implies $\operatorname{rk} E_1^{-t}(K) = 0$.

We now restrict to $K = \mathbb{Q}$ for simplicity.

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$$\begin{aligned} \operatorname{Sel}^2(E/\mathbb{Q}) &\subseteq H^1(G_{\mathbb{Q}}, E[2]) \cong H^1(G_{\mathbb{Q}}, \mathbb{F}_2^2) \cong (\mathbb{Q}^*/\mathbb{Q}^{*2})^2 \\ &= \{(x_1, x_2) : x_1, x_2 \text{ squarefree}\}. \end{aligned}$$

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The local conditions $im(\delta_v)$ are hard to describe for $v \mid \Delta_{E_1}$. But for E_1^d (with *d* squarefree and $gcd(d, \Delta_{E_1}) = 1$), we have the explicit description

$$\operatorname{im}(\delta_{v}) = \begin{cases} \langle (\alpha\beta, d\alpha), (-d\alpha, -\alpha\gamma) \rangle & \text{if } v \mid d \\ H^{1}_{\operatorname{ur}}(G_{\mathbb{Q}_{v}}, \mathbb{F}_{2}^{2}) & \text{if } v \nmid d\Delta_{E_{1}} \end{cases}$$

with $\alpha = a_1 - a_2$, $\beta = a_1 - a_3$ and $\gamma = a_2 - a_3$.

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with $\alpha = a_1 - a_2$, $\beta = a_1 - a_3$ and $\gamma = a_2 - a_3$.

In particular, if $(x_1, x_2) \in \text{Sel}^2(E_1^d/\mathbb{Q})$, then x_1, x_2 are only divisible by primes dividing $d\Delta_E$.

Let E_1/\mathbb{Q} be an elliptic curve with $E_1[2] \cong \mathbb{F}_2^2$. Let $d_1 = p_1 \cdots p_r$ be squarefree and let $d_2 = q_1 \cdots q_r$ be squarefree. Assume that:

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In particular, fixing some *d*, we can compute $\operatorname{Sel}^2(E_1^{dp}/\mathbb{Q})$ from $\operatorname{Sel}^2(E_1^d/\mathbb{Q})$ and the congruence of *p* modulo $8d\Delta_E$.

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Distribution of $\operatorname{Sel}^2(E_1^d/\mathbb{Q})$ was found by Heath-Brown, Kane and Smith.

$$t = m(n - a_1m)(n - a_2m)(n - a_3m) = \kappa P_1 P_2 P_3 P_4$$

with

$$n - a_1 m = P_1, \quad n - a_2 m = P_2, \quad n - a_3 m = P_3, \quad m = \kappa P_4.$$

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(1) First arrange $\operatorname{Sel}^2(E_1^{\kappa}/\mathbb{Q})$ favorably.

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The strategy is then:

- (1) First arrange $\operatorname{Sel}^2(E_1^{\kappa}/\mathbb{Q})$ favorably.
- (2) Then twist by P_1, P_2, P_3, P_4 in a controlled way.
However, since $P_i = n - a_i d$ and $P_4 = \kappa d$, we see that there are constraints between (P_i/P_j) , and moreover $(P_iP_j/r) = +1$ with $r \mid \kappa$.

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We circumvent this by carefully choosing κ and then P_1, P_2, P_3, P_4 (as an analogy: one can choose 3 entries of a 2 × 2-matrix so that the matrix is invertible no matter the choice of the fourth entry).

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It is only here that we use the presence of the 32 real places of K to play off some symbols against each other with quadratic reciprocity.

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Thank you for your attention!