# HOW MANY POTATOES ARE IN A MESH?

Marc van Kreveld Maarten Löffler János Pach

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A LITTLE HISTORY









• Subset P of  $\mathbb{R}^2$ 





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- $\forall p, q \in P : \overline{pq} \subset P$





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- Subset P of  $\mathbb{R}^2$
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- (Also known as "convex set")





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#### "POTATO PEELING PROBLEM"



- Subset P of  $\mathbb{R}^2$
- $\forall p, q \in P : \overline{pq} \subset P$



(Also known as "convex set")



"POTATO PEELING PROBLEM" Find the largest potato contained in a given polygon.

- Subset P of  $\mathbb{R}^2$
- $\forall p, q \in P : \overline{pq} \subset P$



(Also known as "convex set")



"POTATO PEELING PROBLEM" Find the largest potato contained in a given polygon.

(Goodman, 1981)









• Collection of triangles





#### • Collection of triangles





# DEFINITION "MESH"

- Collection of triangles
- Any two triangles:
  - are disjoint,
  - share a single vertex,
  - or share an edge





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• It's a potato ...



• It's a potato ...





- It's a potato ...
- ... and it's also a mesh





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- It's a potato ...
- ... and it's also a mesh
- (Also known as "triangulation")





- It's a potato ...
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#### "MESHED POTATO PEELING PROBLEM"





- It's a potato ...
- ... and it's also a mesh
- (Also known as "triangulation")

# "MESHED POTATO PEELING PROBLEM"



Find the largest meshed potato contained in a given mesh.

- It's a potato ...
- ... and it's also a mesh
- (Also known as "triangulation")

# "MESHED POTATO PEELING PROBLEM"



Find the largest meshed potato contained in a given mesh.

(Aronov et al., 2007)

NEW STUFF













#### <u>DEFINITION</u> "POTATO NUMBER" $\mathcal{P}(M)$





# <u>DEFINITION</u> "POTATO NUMBER" $\mathcal{P}(M)$

 Total number of distinct meshed potatoes in a given mesh M














































































$$\mathcal{P}(\mathsf{M}) = 9$$







• There exists a mesh with potato number  $\Omega(1.5028^{\rm n}).$ 



- There exists a mesh with potato number  $\Omega(1.5028^{n}).$
- For every mesh, the potato number is  $O(1.6181^{n})$ .

# WARNING: TECHNICAL DETAILS

 Place n points in convex position

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- Place n points in convex position
- Triangulate iteratively



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- Place n points in convex position
- Triangulate iteratively
- Count potatoes



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  LEMMA



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**LEMMA** 



The number of potatoes in M grows exponentially with base > 1.5028

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**LEMMA** 



The number of potatoes in M grows exponentially with base > 1.5028 THEOREM

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**LEMMA** 



The number of potatoes in M grows exponentially with base > 1.5028 THEOREM

There is a mesh M with  $\mathcal{P}(M)$  in  $\Omega(1.5028^{n})$ .



 Only count all potatoes containing p



 Only count all potatoes
 containing p



- Only count all potatoes containing p
- Project all vertices from p and remove hit edges



- Only count all potatoes
   containing p
- Project all vertices from p and remove hit edges



- Only count all potatoes containing p
- Project all vertices from p and remove hit edges



- Only count all potatoes containing p
- Project all vertices from p and remove hit edges



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• Orient edges around p

- Only count all potatoes
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• Orient edges around p

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- Orient edges around p
- Call resulting graph G

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#### <u>LEMMA</u>

The number of potatoes in M containing p is bounded by the number of cycles in G



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#### <u>LEMMA</u>

Every edge on the outer face of G either comes from a vertex with outdegree 1, or goes to a vertex with indegree 1.



Let F be a subset of fixed edges of G



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- Let F be a subset of fixed edges of G
- The potential k(G, F) is the #vertices #fixed edges



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 Let Q(k) be the max <sup>K(F)</sup> number of potatoes in a graph with potential k


























- Consider an edge e on the outer face of G
- Then either fix e, or remove its vertices.





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e

 $\mathbf{O}$ 

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## **THEOREM**

For every mesh M,  $\mathcal{P}(M)$  is in  $O(1.6181^{n})$ .

MORE NEW STUFF

• Triangle where all angles are  $> \delta$ 

## • Triangle where all angles are $> \delta$

#### DEFINITION "FAT POTATO"

• Triangle where all angles are  $>\delta$ 

#### DEFINITION "FAT POTATO"



• Triangle where all angles are  $> \delta$ 

## DEFINITION "FAT POTATO"

 Potato that contains a large fat triangle



- Triangle where all angles are  $>\delta$ 

## DEFINITION "FAT POTATO"

 Potato that contains a large fat triangle







## DEFINITION "FAT MESH"

• Mesh where all triangles are fat

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### DEFINITION "FAT MESHED POTATO"



• Mesh where all triangles are fat

## DEFINITION "FAT MESHED POTATO"



## DEFINITION "FAT MESH"

Mesh where all triangles are fat

## DEFINITION "FAT MESHED POTATO"

• Fat potato that is also a fat mesh



## DEFINITION "FAT MESH"

Mesh where all triangles are fat

## DEFINITION "FAT MESHED POTATO"

 Fat potato that is also a fat mesh















#### Potato without internal vertices





#### • Potato without internal vertices



## DEFINITION "CARROT"

- Potato without internal vertices
- (Also known as "outerplanar meshed potato")



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<u>DEFINITION</u> "CARROT NUMBER" C(M)

## DEFINITION "CARROT"

- Potato without internal vertices
- (Also known as "outerplanar meshed potato")



- <u>DEFINITION</u> "CARROT NUMBER" C(M)
- Total number of distinct carrots in a given mesh M
• There exists a fat mesh with fat potato number  $\Omega(n^{\frac{1}{2}\lfloor \frac{2\pi}{\delta} \rfloor})$ .

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# THANK YOU!

