## HOW MANY POTATOES ARE IN A MESH?

Marc van Kreveld
Marten Löffler János Pach

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## A LITTLE HISTORY

DEFINITION "POTATO"

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- Subset P of $\mathbb{R}^{2}$


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## "POTATO PEELING PROBLEM"

Find the largest potato contained in a given polygon.

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## "POTATO PEELING PROBLEM" <br> Find the largest potato contained in a given polygon.

(Goodman, 1981)

DEFINITION "MESH"

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- Collection of triangles


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- Any two triangles:
- are disjoint,
- share a single vertex,
- or share an edge



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- It's a potato ...


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- It's a potato ...
- ... and it's also a mesh


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"MESHED POTATO PEELING PROBLEM"


Find the largest meshed potato contained in a given mesh.

## DEFINITION "MESHED POTATO"

- It's a potato ...
- ... and it's also a mesh
- (Also known as "triangulation")
"MESHED POTATO PEELING PROBLEM"


Find the largest meshed potato contained in a given mesh.
(Aronov et al., 2007)

NEW STUFF

## QUESTION

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How many potatoes can there actually be?

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DEFINITION "POTATO NUMBER" $\mathcal{P}(\mathrm{M})$

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How many potatoes can there actually be?

## DEFINITION "POTATO NUMBER" $\mathcal{P}(\mathrm{M})$

- Total number of distinct meshed potatoes in a given mesh M

EXAMPLE

EXAMPLE


EXAMPLE
$<$


EXAMPLE


EXAMPLE
$\checkmark \Delta \forall$


EXAMPLE



EXAMPLE


EXAMPLE



EXAMPLE
-



EXAMPLE


EXAMPLE


EXAMPLE
$\checkmark \Delta \vee$



$$
\mathcal{P}(\mathrm{M})=9
$$



RESULTS

## RESULTS

- There exists a mesh with potato number $\Omega\left(1.5028^{n}\right)$.


## RESULTS

- There exists a mesh with potato number $\Omega\left(1.5028^{n}\right)$.
- For every mesh, the potato number is $\mathrm{O}\left(1.6181^{\mathrm{n}}\right)$.

WARNING: TECHNICAL DETAILS

LOWER BOUND

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- Place n points in convex position


## LOWER BOUND

- Place n points in



## LOWER BOUND

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- Triangulate iteratively



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- Triangulate iteratively
- Count potatoes



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LEMMA


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LEMMA
The number of potatoes in M
grows exponentially with base > 1.5028

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The number of potatoes in M
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THEOREM

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- Triangulate iteratively
- Count potatoes

LEMMA
The number of potatoes in M
grows exponentially with base > 1.5028
THEOREM
There is a mesh M with $\mathcal{P}(\mathrm{M})$ in $\Omega\left(1.5028^{n}\right)$.

UPPER BOUND

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## UPPER BOUND

- Only count all potatoes containing p



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## UPPER BOUND

- Only count all potatoes containing p
- Project all vertices from $p$ and remove hit edges



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- Orient edges around p


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- Call resulting graph G


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The number of potatoes in M containing p is bounded by the number of cycles in G


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Every edge on the outer face of
G either comes from a vertex with outdegree 1, or goes to a vertex with indegree 1.

## POTENTIAL FUNCTION



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- Let F be a subset of fixed edges of $G$



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- The potential $k(G, F)$ is the \#vertices \#fixed edges



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- Let F be a subset of fixed edges of G
- The potential $k(G, F)$ is the \#vertices \#fixed edges

- Let $\mathrm{Q}(\mathrm{k})$ be the max number of potatoes in a graph with potential k

LEMMA

LEMMA
$Q(k) \leq Q(k-1)+Q(k-2)$

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$\mathrm{Q}(\mathrm{k}) \leq \mathrm{Q}(\mathrm{k}-1)+\mathrm{Q}(\mathrm{k}-2)$

- Consider an edge e on the outer face of $G$

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- Then either fix e, or remove its vertices.

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THEOREM
$\mathrm{Q}(\mathrm{k}) \leq \mathrm{Q}(\mathrm{k}-1)+\mathrm{Q}(\mathrm{k}-2)$

- Consider an edge e on the outer face of $G$


- Then either fix e, or remove its vertices.

THEOREM
For every mesh $\mathrm{M}, \boldsymbol{P}(\mathrm{M})$ is in $\mathrm{O}\left(1.6181^{n}\right)$.

MORE NEW STUFF

DEFINITION "FAT TRIANGLE"

## DEFINITION "FAT TRIANGLE"

- Triangle where all angles are $>\delta$


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## DEFINITION "FAT POTATO"

- Potato that contains a large fat triangle


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DEFINITION "FAT MESH"

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- Mesh where all triangles are fat


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DEFINITION "FAT MESHED POTATO"

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DEFINITION "FAT MESHED POTATO"

## DEFINITION "FAT MESH"

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DEFINITION "FAT MESHED POTATO"

- Fat potato that is also a fat mesh


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DEFINITION "CARROT"

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- Potato without internal vertices


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- (Also known as "outerplanar meshed potato")



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DEFINITION "CARROT NUMBER" $\mathcal{C}(\mathrm{M})$

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- (Also known as "outerplanar meshed potato")



## DEFINITION "CARROT NUMBER" $\mathcal{C}(\mathrm{M})$

- Total number of distinct carrots in a given mesh M

MORE RESULTS

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- There exists a fat mesh with fat potato number $\Omega\left(n^{\frac{1}{2}\left\lfloor\frac{2 \pi}{\delta}\right\rfloor}\right)$.


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- There exists a fat mesh with fat potato number $\Omega\left(n^{\frac{1}{2}\left\lfloor\frac{2 \pi}{\sigma}\right\rfloor}\right)$.
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## MORE RESULTS

- There exists a fat mesh with fat potato number $\Omega\left(n^{\frac{1}{2}\left\lfloor\frac{2 \pi}{\delta}\right\rfloor}\right)$.
- For every fat mesh, the potato number is $\mathrm{O}\left(\mathrm{n}^{\left\lceil\frac{\pi}{8}\right\rceil}\right)$.

- There exists a fat mesh with fat carrot number $\Omega\left(\mathrm{n}^{\left\lfloor\frac{2 \pi}{35}\right\rfloor}\right)$.


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- For every fat mesh, the potato number is $\mathrm{O}\left(\mathrm{n}^{\left\lceil\frac{\pi}{8}\right\rceil}\right)$.

- There exists a fat mesh with fat carrot number $\Omega\left(\mathrm{n}^{\left\lfloor\frac{2 \pi}{36}\right\rfloor}\right)$.



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- There exists a fat mesh with fat potato number $\Omega\left(n^{\frac{1}{2}\left\lfloor\frac{2 \pi}{\delta}\right\rfloor}\right)$.
- For every fat mesh, the potato number is $\mathrm{O}\left(\mathrm{n}^{\left\lceil\frac{\pi}{8}\right\rceil}\right)$.

- There exists a fat mesh with fat carrot number $\Omega\left(n^{\left\lfloor\frac{2 \pi}{36}\right\rfloor}\right)$.
- For every fat mesh, the carrot number is $\left.\left.\mathrm{O}\left(\mathrm{n}^{2 \pi}\right\rfloor \mathrm{L}\right\rfloor\right)$.



## THANK YOU!



