

REGION-BASED APPROXIMATION ALGORITHMS FOR VISIBILITY BETWEEN IMPRECISE LOCATIONS

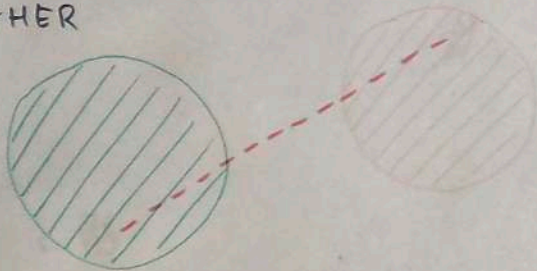
KEVIN BUCHIN • IRINA KOSTITSYNA • MAARTEN LÖFFLER • RODRIGO S.

ACT I

REGION-BASED APPROXIMATION WHAT?

PROBLEM STATEMENT

COMPUTE THE PROBABILITY THAT TWO RANDOM POINTS CAN SEE EACH OTHER



MONKEY BEHAVIOUR

DO WANDERING MONKEY GROUPS ALTER THEIR COURSE WHEN THEY SEE EACH OTHER?



CONVEXITY MEASURE

CAN WE PARAMETERISE THE EXTENT TO WHICH A POLYGON IS CONVEX?

[ERIK WILLEMS, 2013]

BETWEEN THREE

NA KOSTITSYNA • MAARTEN LÖF

ACT I

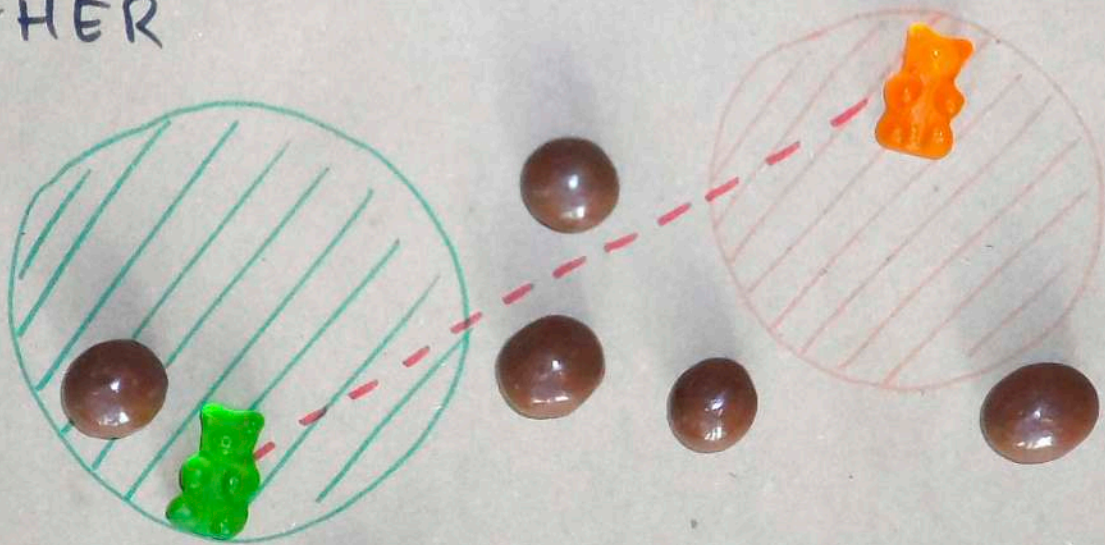
REGION-BASED APPROXIMATION WHAT?

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LITY THAT

PROBLEM STATEMENT

COMPUTE THE PROBABILITY THAT
TWO RANDOM POINTS CAN SEE EACH
OTHER



CONVEXITY M

CAN WE PARAMETERISE

MONKEY BEHAVIOUR

DO WANDERING MONKEY GROUPS
ALTER THEIR COURSE WHEN THEY
SEE EACH OTHER?

ASURE

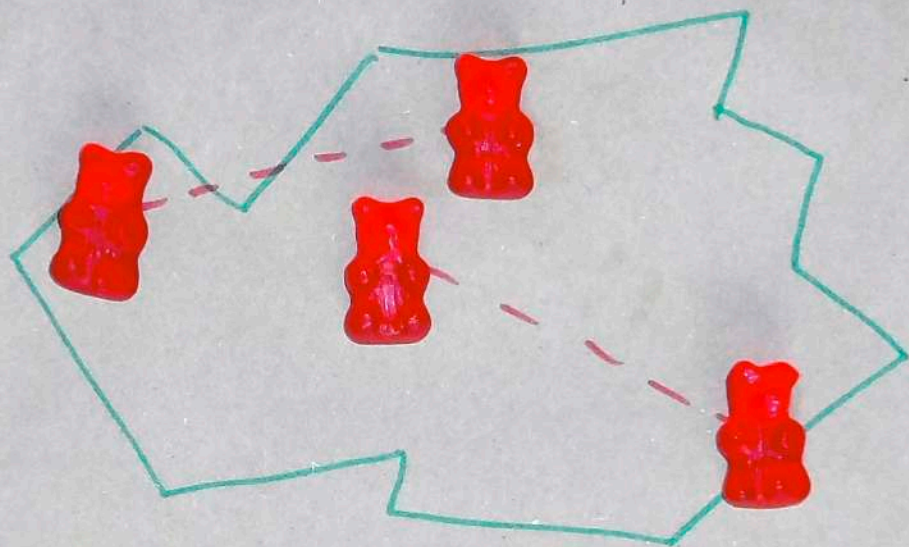
THE EXTENT
CONVEX?



[ERIK WILLEMS, 2013]

CONVEXITY MEASURE

CAN WE PARAMETERISE THE EXTENT
TO WHICH A POLYGON IS CONVEX?



[GÜNTERRÖTE, 2012]



[GÜNTERRÖTE, 2012]

ACT II

WAIT, ISN'T THAT A SOLVED PROBLEM?

RATION

CK YOUR SEAT
ANDOM HATS

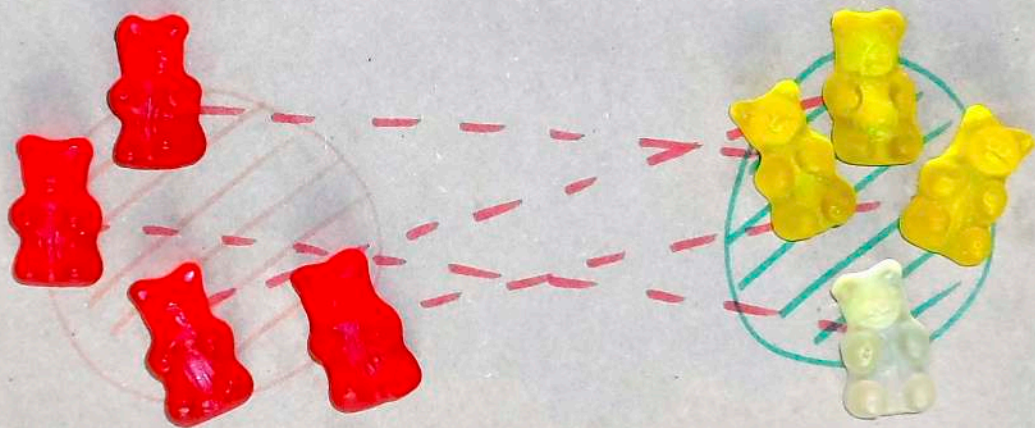
POINT SA

JUST TRY TWO
AND TEST THEIR



POINT SAMPLING

JUST TRY TWO RANDOM POINTS,
AND TEST THEIR VISIBILITY!



(AND THEN DO IT AGAIN. AND AGAIN.)

CKS

RELIABLE SOURCE

ILLUSTRATION

PLEASE CHECK YOUR SEAT
FOR ANY RANDOM HATS



D
- Y
O
- Y
S
- M

(AND THEN D

DRAWBACKS

- YOU NEED A RELIABLE SOURCE OF RANDOMNESS
- YOU NEED TO TAKE A LOT OF SAMPLES
- MOST LIKELY YOU STILL WON'T GET THE RIGHT ANSWER
- YOU JUST NEVER KNOW FOR SURE!

- YOU NEED TO TAKE A LOT OF SAMPLES
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ACT III

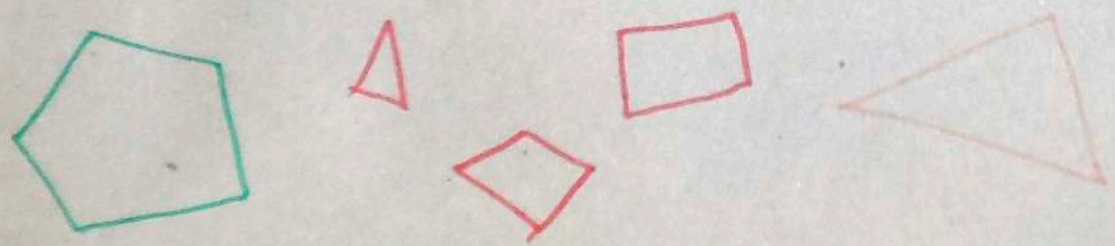
A FUNDAMENTAL APPROACH,
WITH INTEGRALS AND STUFF

BLEM

LYGONS OF
COLLECTION

FORMAL PROBLEM

GIVEN TWO CONVEX POLYGONS OF COMPLEXITY N AND A COLLECTION OF OBSTACLES OF COMPLEXITY M , COMPUTE THE PROBABILITY THAT TWO POINTS, TAKEN FROM BOTH POLYGONS, ARE MUTUALLY VISIBLE



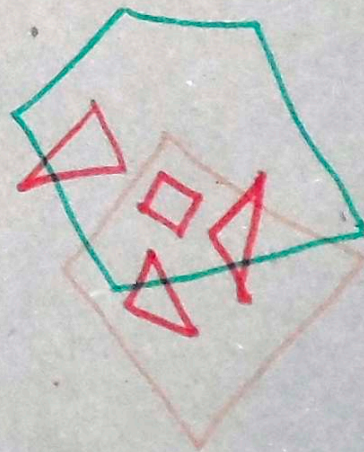
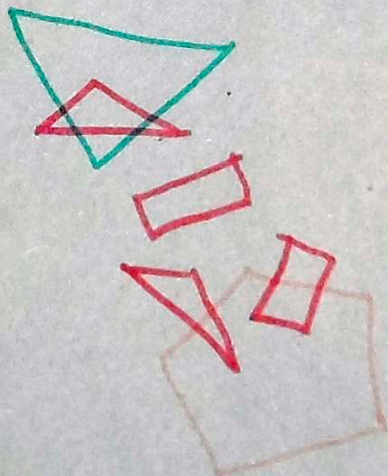
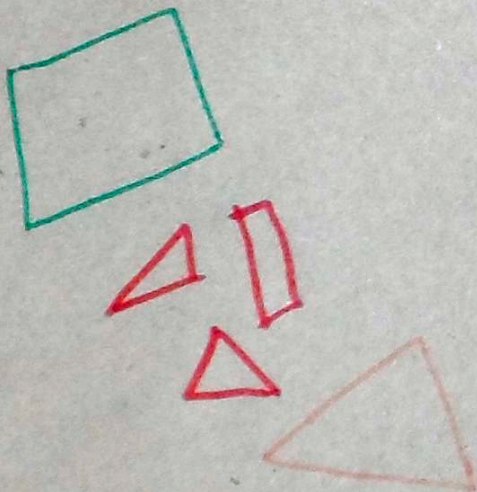
SOME
OT



THAT
BOTH
SIBLE

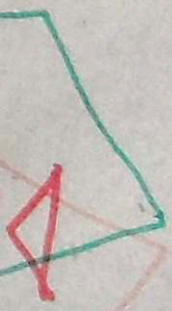
VARIANTS

SOME POLYGONS OVERLAP.
OTHERS DON'T.



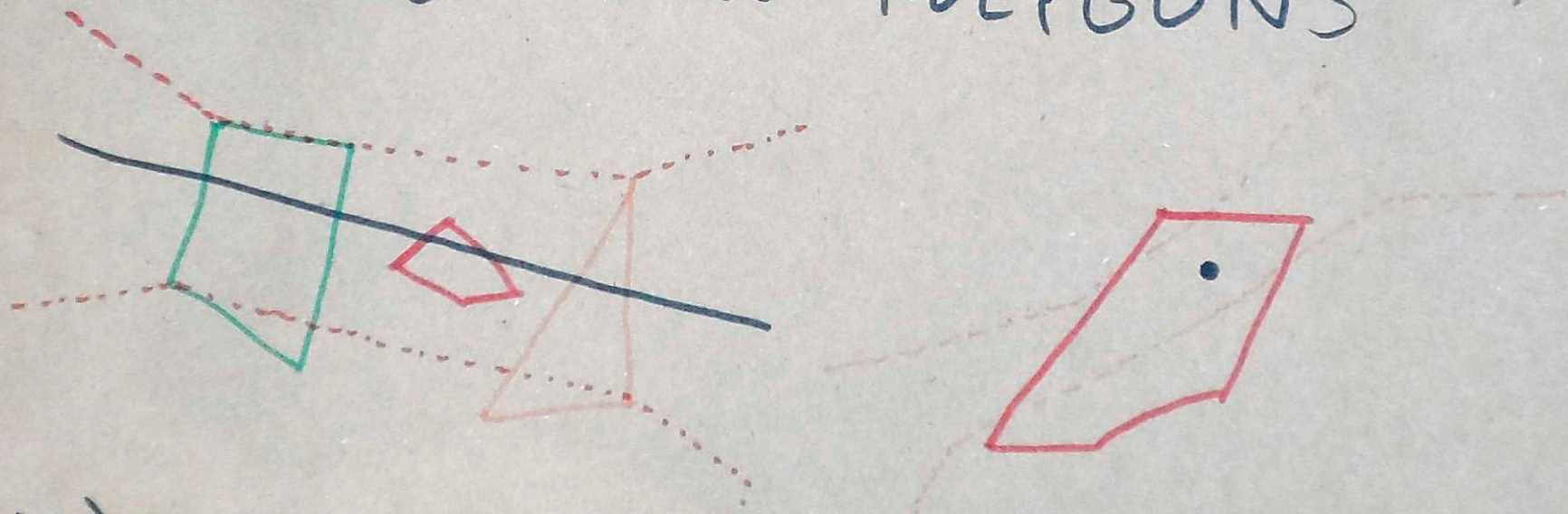
POI

WE W
THE S



POINT-LINE DUALITY

WE WANT TO ARGUE ABOUT
THE SPACE OF ALL LINES
THROUGH TWO POLYGONS



THE CALCULATION

$$I = \iint_C \left(\int_{X_1(\alpha, \beta)}^{X_2(\alpha, \beta)} \int_{X_3(\alpha, \beta)}^{X_4(\alpha, \beta)} (x_2 - x_1) dx_2 dx_1 \right) d\alpha d\beta = \iint_C \frac{(x_2 - x_1)^2}{2} d\alpha d\beta$$

$$I_{ij} = \iint_C X_i X_j^2 d\alpha d\beta = \iint_C \frac{(b_i \beta - c_i)(b_j \beta - c_j)^2}{(b_i \alpha + a_i)(b_j \alpha + a_j)^2} d\alpha d\beta = \sum_{i,j} \dots$$

$$F_{ij}(\alpha) = \int_{A_1 \alpha + B_1}^{A_2 \alpha + B_2} \frac{(b_i \beta - c_i)(b_j \beta - c_j)^2}{(b_i \alpha + a_i)(b_j \alpha + a_j)^2} d\beta = \frac{\log(a_i + \alpha b_i)}{12 b_i^2 b_j^2} \left[4(A_2 \alpha + B_2) \right. \\ \left. + \frac{1}{24 b_i^2 b_j^2} \left[3\alpha^2 (a_i + \alpha b_i) \right] \right]$$

$$= \iint_C \frac{(x_2 - x_1)(x_4 - x_3)(x_3 + x_4 - x_1 - x_2)}{2} dx d\beta = \frac{1}{2} \iint_C (-x_1^2 + x_2^2 + x_3^2 + x_4^2) dx d\beta$$

$$- dx d\beta = \sum_{C \vee C} \int_{\alpha_1}^{\alpha_2} \left(\int_{A_1 \alpha + B_1}^{A_2 \alpha + B_2} \frac{(b_i \beta - c_i)(b_j \beta - c_j)^2}{(b_i \alpha + a_i)(b_j \alpha + a_j)^2} d\alpha \right) d\beta$$

$$\frac{+ \alpha b_i}{b_i^2} \left[4(A_2^3 - A_1^3)(a_j b_i - a_i b_j)(b_i c_j - b_j c_i) + 3(A_2^4 - A_1^4) \right]$$

$$\frac{2 b_j^2}{b_j^2} \left[3 \alpha^2 (A_2^4 - A_1^4) b_i b_j + \alpha (A_2^4 - A_1^4) (12 a_j b_i - 6 a_i b_j) \right]$$

CALCULATION

$$\int_C \left(\int_{x_2(\alpha, \beta)}^{x_3(\alpha, \beta)} \int_{x_1(\alpha, \beta)}^{x_2(\alpha, \beta)} (x_2 - x_1) dx_1 dx_2 \right) d\alpha d\beta = \int_C$$

$$I_{ij} = \iint_C x_i x_j^2, d\alpha d\beta = \iint_C \frac{(b_i \beta - c_i)(b_j \beta - c_j)^2}{(b_i \alpha + a_i)(b_j \alpha + a_j)^2} d\alpha d\beta = \int_C$$

$$F_{ij}(\alpha) = \int_{A_i \alpha + B_i}^{A_j \alpha + B_j} \frac{(b_j \beta - c_j)(b_j \beta - c_j)^2}{(b_i \alpha + a_i)(b_j \alpha + a_j)^2} d\beta = \frac{\log \frac{(a_j + b_j)}{(a_i + b_i)}}{12 b_j^2 c_j^2} \left[\frac{4(A_j^3 - A_i^3)}{12 b_j^2 c_j^2} \right]$$

$$+ \frac{1}{24 b_j^2 c_j^2} \left[3x^2(A_j - A_i) \right]$$

OTHER CASES ARE SIMILAR BUT SLIGHTLY MORE COMPLICATED. PLEASE THEOREM THE DESIRED PROBABILITIES CAN BE COMPUTED IN THE

THEOREM

THE DESIRED PROBABILITY
CAN BE COMPUTED IN TIME

$$O(M \cdot (M+N)^2)$$

ACT IV

LIGHTLY MORE COMPUTATION

THEOREM

THE DESIRED PROBABILITY
CAN BE COMPUTED IN TIME

$$O(M \cdot (M+N)^2)$$

ACT IV

SURELY WE CAN DO BETTER?

ANCE

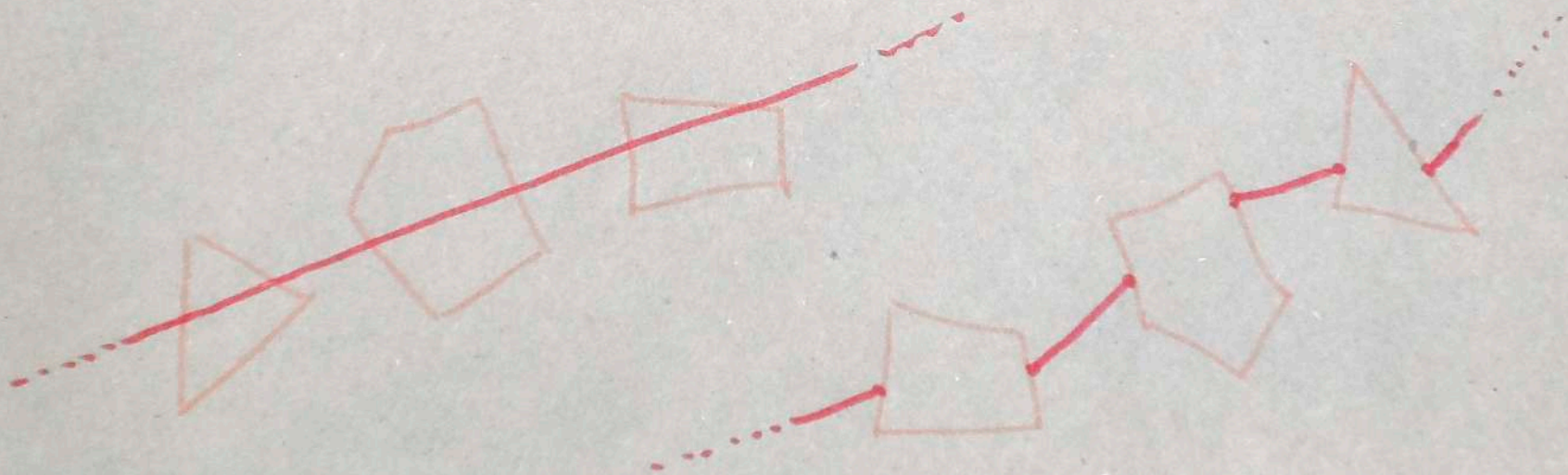
OF
END

LAYER

EACH LINE
MAXIMAL SEC

SEGMENT SPACE

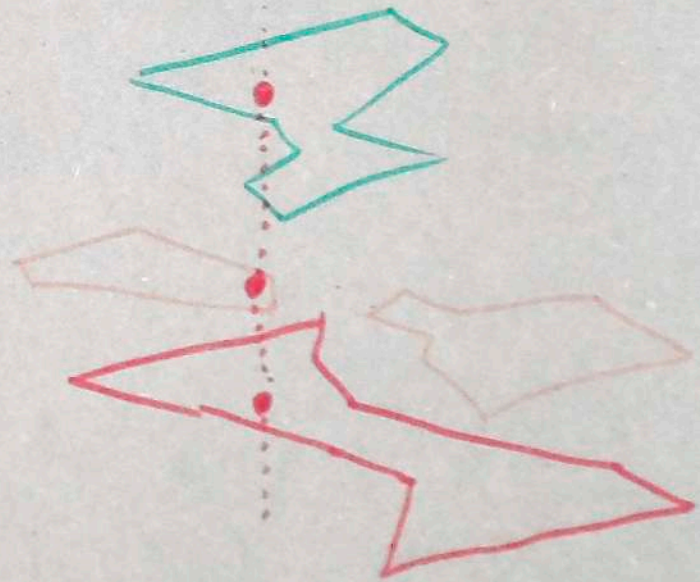
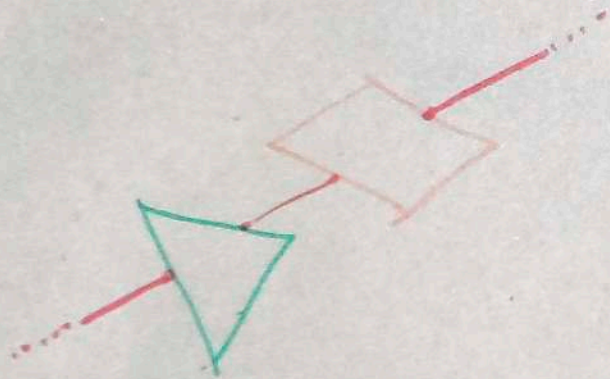
CONSIDER THE SPACE OF
LINE SEGMENTS WHOSE END
POINTS TOUCH AN OBSTACLE



TH

LAYERS OF SEGMENTS

EACH LINE MAY CONSIST OF M
MAXIMAL SEGMENTS, ADDING A DISCRETE
DIMENSION TO THE SEGMENT SPACE



THEOREM

THE DESIRED PROBABILITY
CAN ALSO BE COMPUTED IN

JUST $O((M+N)^2)$ TIME

THE DESIRED PROBABILITY
CAN ALSO BE COMPUTED IN
JUST $O((M+N)^2)$ TIME

ACT V

BUT WHAT ABOUT REAL DISTRIBUTIONS?

DON'T

GAUSSIANS

REAL RANDOM POINTS DON'T
HAVE UNIFORM DISTRIBUTIONS



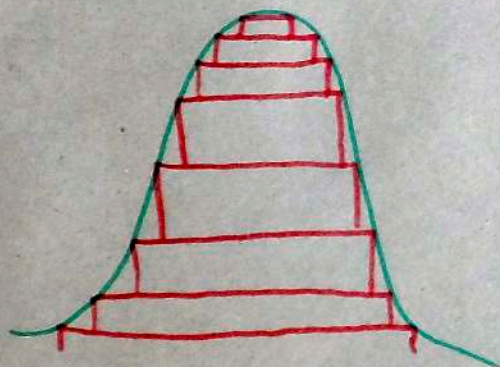
FOR NORMAL DISTRIBUTIONS,
OUR INTEGRAL HAS NO CLOSED-
FORM SOLUTION!

AP

WE
DIS
DIS

APPROXIMATION BY DISKS

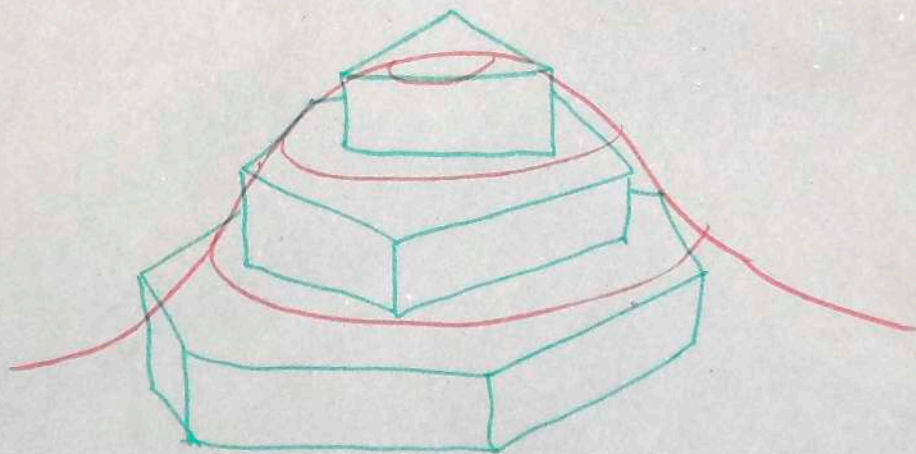
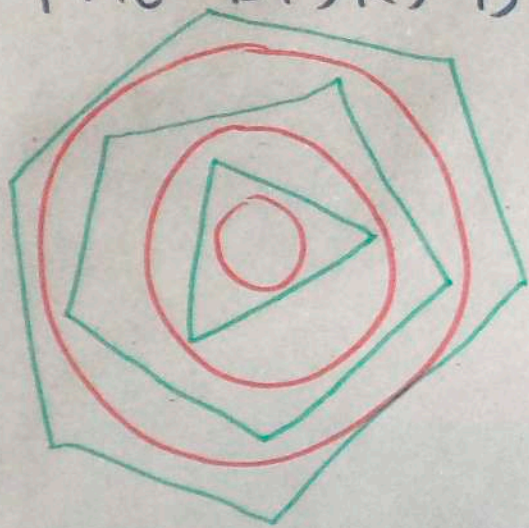
WE CAN APPROXIMATE A GAUSSIAN
DISTRIBUTION WITH A BUNCH OF UNIFORM
DISTRIBUTIONS ON CONCENTRIC DISKS



WITH POLYGONS

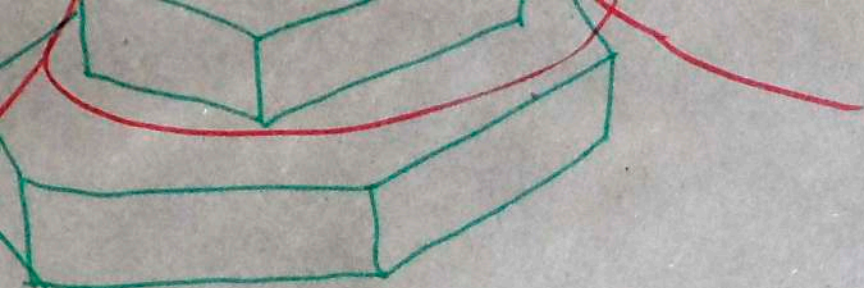
APPROXIMATION WITH POLYGONS

AND THEN, WE CAN IN TURN APPROXIMATE
THE DISKS BY REGULAR POLYGONS



THEOREM

A NORMAL DISTRIBUTION
CAN BE ε -APPROXIMATED
WITH POLYGONS OF TOTAL
COMPLEXITY $O(\sqrt{\varepsilon})$



WITH POLYGON
COMPLEXITY

ACT VI

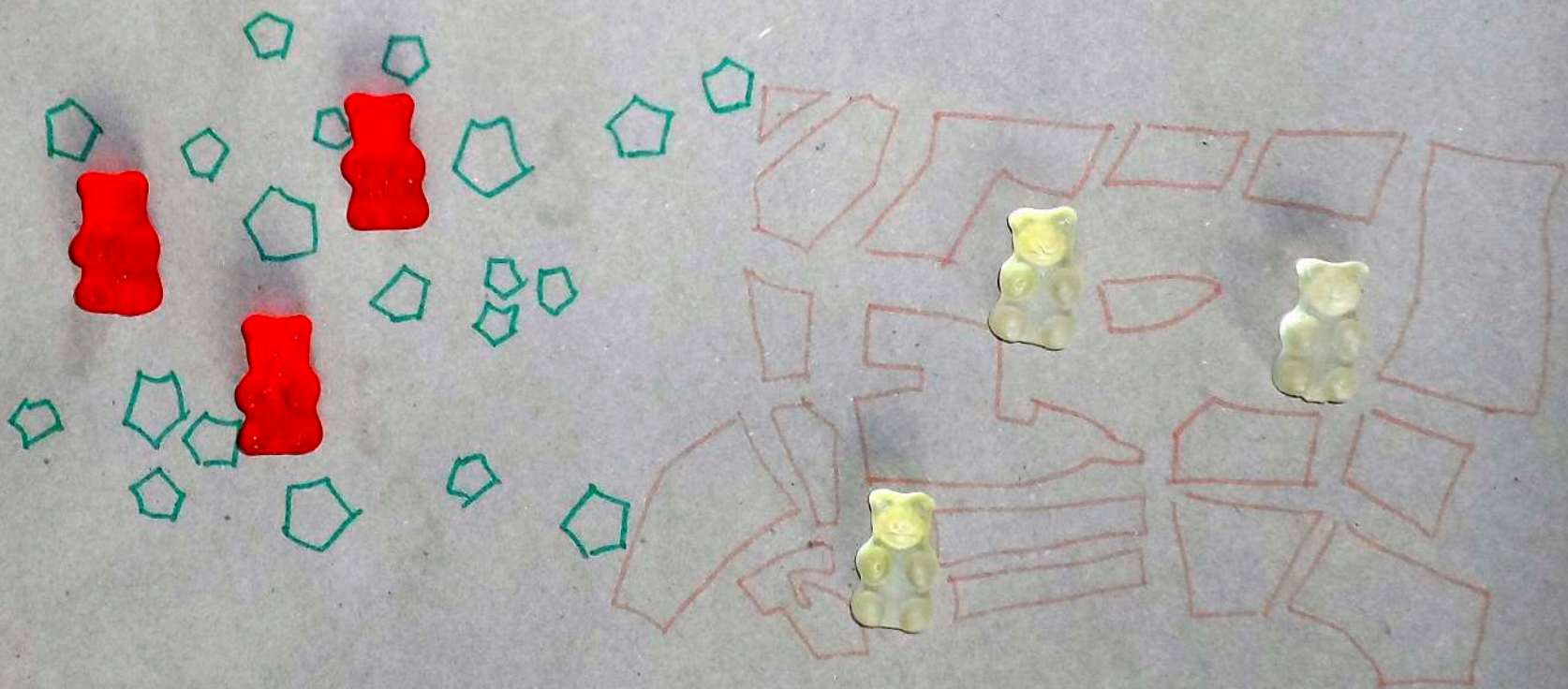
ISN'T THIS SUPPOSED TO
BE AN EXPERIMENTAL PAPER?

SET

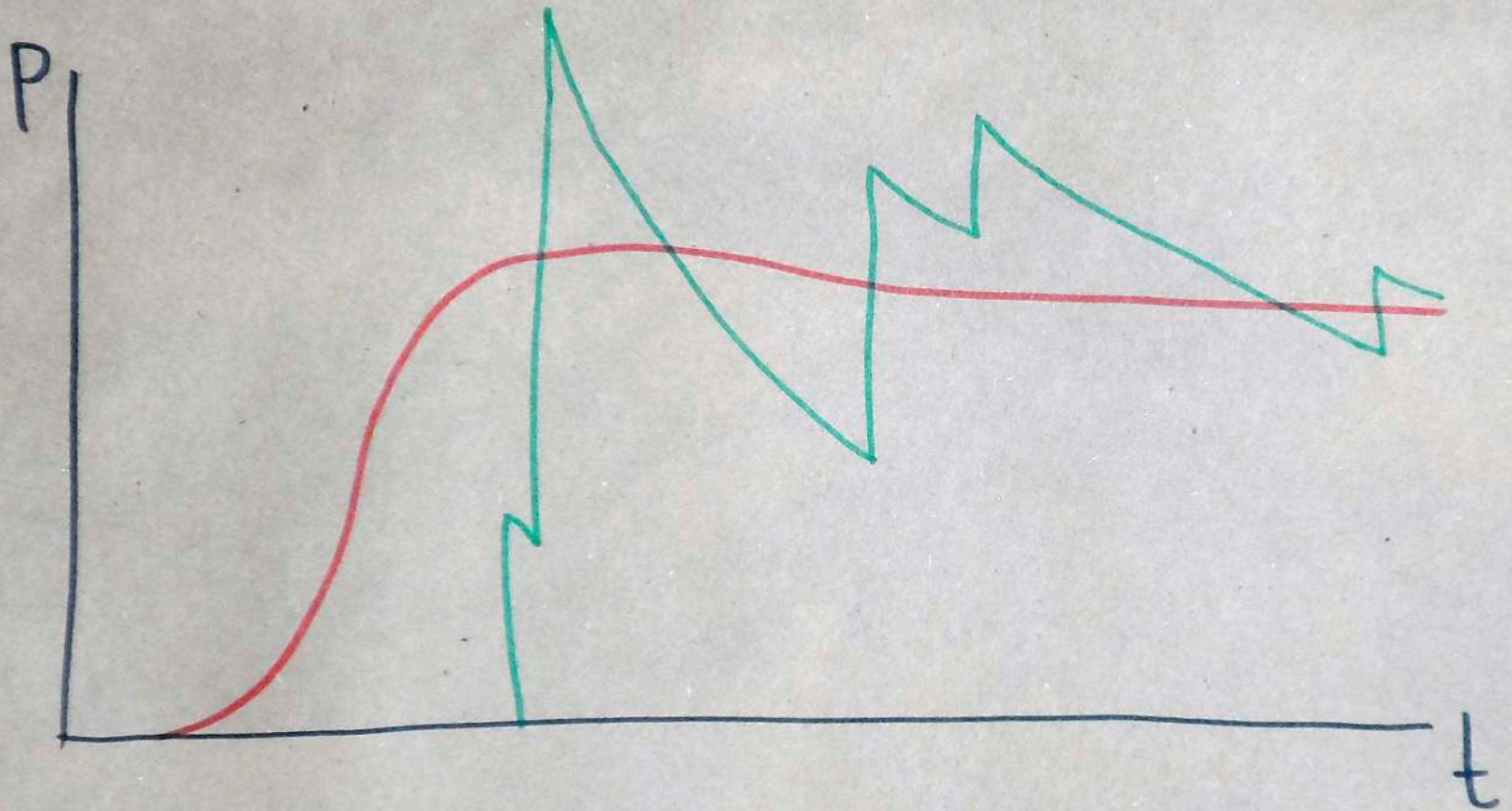
WE PER
A FORE

SETTINGS

WE PERFORMED TESTS IN
A FOREST AN URBAN ENVIRONMENT



RESULTS



— t

THE END