Seminar on Set Theory

Solutions to exercise 2

September 25, 2015

## 1 Boolean Algebras and Propositional Logic (6 points total)

Let  $\mathcal{L}$  be the classical propositional language consisting of  $\bot, \top, \neg, \rightarrow, \lor, \land, \leftrightarrow$ and propositional variables  $P_x$  for each  $x \in B$ . Let the theory T have the following sentences for all  $x, y \in B$ :

$$P_x \wedge P_y \leftrightarrow P_{x \wedge y},\tag{1}$$

$$P_{x^*} \leftrightarrow \neg P_x,\tag{2}$$

$$\neg(P_x \to P_y) \quad \text{if } x \not\leq y.$$
 (3)

(Students get 1,5 points for defining a good theory)

Define a map  $f: B \to B(T)$  by  $x \mapsto [P_x]$  for all  $x \in B$ . We will show that f is a bijective algebra homomorphism.

If  $x, y \in B$ , then

$$f(x \wedge y) = [P_{x \wedge y}]$$
  
=  $[P_x \wedge P_y]$  (by (1))  
=  $[P_x] \wedge [P_y]$  (by construction of  $B(T)$ )  
=  $f(x) \wedge f(y)$ .

Furthermore,

$$f(x^*) = [P_{x^*}]$$
  
=  $[\neg P_x]$  (by (2))  
=  $[P_x \rightarrow \bot]$   
=  $[P_x] \rightarrow [\bot]$  (by construction of  $B(T)$ )  
=  $[P_x]^*$   
=  $f(x)^*$ .

Proposition 1.1 now tells us that f is an algebra homomorphism. (Students get 1 point for showing this)

To show that f is injective, suppose we have  $x, y \in B$  with  $x \neq y$ . We may assume that  $x \not\leq y$ . By (3), we must have that  $T \vdash \neg (P_x \rightarrow P_y)$ , so

 $T \not\vdash P_x \to P_y$ . Hence, by construction of B(T), we have  $f(x) = [P_x] \neq [P_y] = f(y)$ . (Students get 2 point for showing injectivity)

We prove that f is surjective by induction on the complexity of formulas of B(T). Clearly, for every propositional variable  $P_x$  of  $\mathcal{L}$ , there is an  $x \in$ B such that  $f(x) = [P_x]$ . Furthermore,  $f(0) = [\bot]$  and  $f(1) = [\top]$  (by Proposition 1.1). Assume that there are  $x, y \in B$  such that  $f(x) = [\phi]$  and  $f(y) = [\psi]$ , then clearly:

$$\begin{split} f(x^*) &= [\phi]^* = [\neg \phi], \\ f(x \wedge y) &= [\phi] \wedge [\psi] = [\phi \wedge \psi], \\ f(x \vee y) &= [\phi] \vee [\psi] = [\phi \vee \psi], \\ f(x \Rightarrow y) &= [\phi] \Rightarrow [\psi] = [\phi \rightarrow \psi], \\ f((x \Rightarrow y) \wedge (y \Rightarrow x)) &= [\phi \leftrightarrow \psi] \quad \text{(by the equalities above).} \end{split}$$

(Students get 2 points for showing surjectivity) We conclude that f is a bijective algebra homomorphism, so B and B(T) are isomorphic.

## 2 Cantor's Theorem (4 points total)

a) Suppose Cantor's Theorem is not true. Then there is a set X and a bijection  $f: X \to \mathcal{P}(X)$  (It is trivial to show that the power set of a set is not of lower cardinality than that set). Clearly X is nonempty because  $\emptyset$  and  $\{\emptyset\}$  have different finite cardinality. Define the subset  $X_0 \subseteq X$  by:

$$X_0 = \{ x \in X : x \notin f(x) \}.$$

By Zermelo's third Axiom of separation this indeed is a set. And it is a subset of X, so  $X_0 \in \mathcal{P}(X)$ . Since f is a bijection, there must be a  $x \in X$  such that  $X_0 = f(x)$ . Like in Zermelo's proof we again consider the two possible cases:

 $x \in X_0$ : Then, by definition of  $X_0, x \notin f(x) = X_0$  which is a contradiction.

 $x \in X \setminus X_0$ : Then x is an element of X, not in f(x). So  $x \in X_0$  which again is a contradiction.

Thus, we get a contradiction and conclude that there is no such bijection. Thereby proving the theorem. (Students get 3 points for Proving the theorem)

b) Suppose such a set U does exist. Then by the powerset axiom  $\mathcal{P}(U)$  is also a set and all its elements are contained in U. So  $|\mathcal{P}(U)| \leq |U|$  which is in contradiction with Cantor's theorem. (Students get 1 point for showing this)