# Seminar on Set Theory

Model solution exercise 8

December 1, 2015

## 3 Some details about an action on ${\cal V}^{(B)}$

In this exercise you will prove some omitted details from the proof of Theorem 3.3. Define the map  $\langle g, u \rangle \mapsto gu : G \times V^{(B)} \to V^{(B)}$  by recursion on the well-founded relation  $y \in \operatorname{dom}(x)$  via

$$gu = \{ \langle gx, g \cdot u(x) \rangle : x \in \operatorname{dom}(u) \}.$$

a)

Prove that

 $g \cdot \llbracket u \in v \rrbracket = \llbracket gu \in gv \rrbracket$ 

and

$$g \cdot \llbracket u = v \rrbracket = \llbracket gu = gv \rrbracket.$$

(2 pt.)

#### Solution:

Use induction on the well-founded relation defined on page 23. Assume that  $g \cdot [\![x \in y]\!] = [\![gx \in gy]\!]$  and  $g \cdot [\![x = y]\!] = [\![gx = gy]\!]$  for all  $\langle x, y \rangle < \langle u, v \rangle$ . Then

$$\llbracket gu \in gv \rrbracket = \bigvee_{x \in \operatorname{dom}(gv)} (gv(x) \land \llbracket gu = x \rrbracket)$$
$$= \bigvee_{y \in \operatorname{dom}(v)} (gv(gy) \land \llbracket gu = gy \rrbracket)$$
(by Thm 3.3 (i) and IH)
$$= \bigvee_{y \in \operatorname{dom}(v)} (g \cdot v(y) \land g \cdot \llbracket u = y \rrbracket)$$
$$= g \cdot \bigvee_{y \in \operatorname{dom}(v)} (v(y) \land \llbracket u = y \rrbracket)$$
$$= g \cdot \llbracket u \in v \rrbracket.$$

Also we have that

$$\begin{split} \llbracket gu = gv \rrbracket &= \bigwedge_{x' \in \operatorname{dom}(gu)} (gu(x') \Rightarrow \llbracket x' \in gv \rrbracket) \land \bigwedge_{y' \in \operatorname{dom}(gv)} (gv(y') \Rightarrow \llbracket y' \in gu \rrbracket) \\ &= \bigwedge_{x \in \operatorname{dom}(u)} (gu(gx) \Rightarrow \llbracket gx \in gv \rrbracket) \land \bigwedge_{y \in \operatorname{dom}(v)} (gv(gy) \Rightarrow \llbracket gy \in gu \rrbracket) \\ (\text{by Thm 3.3 (i) and IH}) &= \bigwedge_{x \in \operatorname{dom}(u)} (g \cdot u(x) \Rightarrow g \cdot \llbracket x \in v \rrbracket) \land \bigwedge_{y \in \operatorname{dom}(v)} (g \cdot v(y) \Rightarrow g \cdot \llbracket y \in u \rrbracket) \\ &= g \cdot \Big(\bigwedge_{x \in \operatorname{dom}(u)} (u(x) \Rightarrow \llbracket x \in v \rrbracket) \land \bigwedge_{y \in \operatorname{dom}(v)} (v(y) \Rightarrow \llbracket y \in u \rrbracket) \Big) \\ &= g \cdot \Big[ u = v \rrbracket. \end{split}$$

### b)

Prove Theorem 3.3 (ii):  $g\hat{v} = \hat{v}$  for any  $v \in V$ . (3 pt.)

### Solution:

By induction on the well-founded relation  $\in.$  Suppose  $g\hat{y}=\hat{y}$  for all  $y\in v.$  Then

$$\begin{split} g\hat{v} &= \{\langle gx, g \cdot \hat{v}(x) \rangle | x \in \operatorname{dom}(\hat{v}) \} \\ &= \{\langle g\hat{y}, g \cdot \hat{v}(\hat{y}) \rangle | y \in v \} \\ &= \{\langle g\hat{y}, g \cdot 1 \rangle | y \in v \} \\ &= \{\langle \hat{y}, 1 \rangle | y \in v \} \\ &= \hat{v}. \end{split}$$