# Seminar on Set Theory

### Model solution 9

## Exercise 1

### Exercise 1a

(a) It suffices to show that  $g(x \wedge y) = g(x) \wedge g(y)$  and  $g(x^*) = g(x)^*$  for all  $x, y \in B'$ . So let  $x, y \in B$  be arbitrary. Since h is a bijection, we have

$$g(x \wedge y) = g(x) \wedge g(y) \longleftrightarrow h(g(x \wedge y)) = h(g(x) \wedge g(y))$$

Now  $h(g(x \land y)) = x \land y$  and  $h(g(x) \land g(y)) = h(g(x)) \land h(g(y)) = x \land y$ , so indeed  $g(x \land y) = g(x) \land g(y)$ . Similarly

$$g(x^*) = g(x)^* \longleftrightarrow h(g(x^*)) = h(g(x)^*).$$

So since  $h(g(x^*)) = x^*$  and  $h(g(x)^*) = h(g(x))^* = x^*$ , we conclude that  $g(x^*) = g(x)^*$ . Hence, g is a homomorphism. (1.5 points)

## Exercise 1b

Note that  $\pi$  and  $\pi^{-1}$  are both order preserving, as they are homomorphisms. Let  $X \subseteq B$  and suppose that  $\bigvee X$  exists in B. Since  $\bigvee X \ge x$  for all  $x \in X$ , we have  $\pi(\bigvee X) \ge \pi(x)$  for all  $x \in X$ , so  $\pi(\bigvee X)$  is an upper bound for  $\{\pi(x) \mid x \in X\}$ . Suppose that  $y \in B$  is also an upper bound for  $\{\pi(x) \mid x \in X\}$ . Then  $y \ge \pi(x)$  for all  $x \in X$ , so  $\pi^{-1}(y) \ge x$  for all  $x \in X$ . It follows that  $\pi^{-1}(y)$  is an upper bound for X, so  $\pi^{-1}(y) \ge \bigvee X$ . This implies  $y \ge \pi(\bigvee X)$ . We conclude that  $\pi(\bigvee X)$  is the least upper bound for  $\{\pi(x) \mid x \in X\}$ , so  $\pi(\bigvee X) = \bigvee \{\pi(x) \mid x \in X\}$ . Hence,  $\pi$  is a complete homomorphism. (1.5 points)

#### Exercise 1c

Suppose that B is homogeneous and let  $x \neq 0, y \neq 0$  be in B. If we let  $\pi' \in Aut(B)$ , then

$$\pi'\left(\bigvee\{\pi(x)\mid \pi\in \operatorname{Aut}(B)\}\right) = \bigvee\{\pi'(\pi(x))\mid \pi\in \operatorname{Aut}(B)\}$$

by part **b**. Furthermore

$$\bigvee \{ \pi'(\pi(x)) \mid \pi \in \operatorname{Aut}(B) \} = \bigvee \{ \pi(x) \mid \pi \in \operatorname{Aut}(B) \}$$

since  $\pi'\pi$  runs through Aut(B) as  $\pi$  runs through Aut(B). This means that  $\bigvee \{\pi(x) \mid \pi \in \text{Aut}(B)\}$  is invariant, so it must have value 0 or 1, by homogeneity of B. Since  $\bigvee \{\pi(x) \mid \pi \in \text{Aut}(B)\} \ge \text{id}(x) = x$  and  $x \ne 0$ , it follows that  $\bigvee \{\pi(x) \mid \pi \in \text{Aut}(B)\} = 1$ . Hence

$$y = y \land \bigvee \{ \pi(x) \mid \pi \in \operatorname{Aut}(B) \} = \bigvee \{ y \land \pi(x) \mid \pi \in \operatorname{Aut}(B) \}.$$

So since  $y \neq 0$ , there must be  $\pi \in \operatorname{Aut}(B)$  such that  $y \wedge \pi(x) \neq 0$ .

Conversely, suppose that B is not homogeneous. Then there exists an invariant element  $y \in B$ , with  $y \neq 0$  and  $y \neq 1$ . But then  $y^* \neq 0$ , so if we take  $x = y^*$ , then we have found nonzero  $x, y \in B$  such that

$$x \wedge \pi(y) = y^* \wedge \pi(y) = y^* \wedge y = 0,$$

for all  $\pi \in Aut(B)$ . (2 points)

# Exercise 2

### Exercise 2a

This is shown by proving  $V_{\alpha}^{(\Gamma)} \subseteq V_{\alpha}^{(B)}$  for any ordinal  $\alpha$ . For  $\alpha$  an ordinal, we can show  $V_{\alpha}^{(\Gamma)} \subseteq V_{\alpha}^{(B)}$  by induction. Assume for all  $\beta < \alpha$  we know  $V_{\beta}^{(\Gamma)} \subseteq V_{\beta}^{(B)}$ . Now let  $x \in V_{\beta}^{(\Gamma)}$ , then  $\operatorname{Fun}(x) \wedge \operatorname{ran}(x) \subseteq B \wedge \exists \beta < \alpha \operatorname{dom}(x) \subseteq V_{\beta}^{(\Gamma)}$ , so by the induction hypothesis  $\operatorname{Fun}(x) \wedge \operatorname{ran}(x) \subseteq B \wedge \exists \beta < \alpha \operatorname{dom}(x) \subseteq V_{\beta}^{(B)}$ , so  $x \in V_{\alpha}^{(B)}$ . So by induction  $V_{\alpha}^{(\Gamma)} \subseteq V_{\alpha}^{(B)}$  for any ordinal  $\alpha$ , so  $V^{(B)} \subseteq V^{(\Gamma)}$ . (.5 points for a correct answer)

### Exercise 2b

Let  $u = \{\langle \emptyset, r \rangle\} \in V^{(B)}$ . Now for  $g \in \operatorname{stab}(u)$  we know  $\{\langle \emptyset, r \rangle\} = u = gu = \{\langle g\emptyset, gr \rangle\} = \{\langle \emptyset, gr \rangle\}$ , i.e. gr = r, so  $\operatorname{stab}(u) \subseteq \operatorname{stab}(r)$ . But as  $\operatorname{stab}(r) \notin \Gamma$ ,  $\operatorname{stab}(u) \notin \Gamma$ (because  $\Gamma$  is a filter of subgroups), so  $u \notin V^{(\Gamma)}$ . So  $V^{(\Gamma)} \neq V^{(B)}$ .(.5 points for a correct answer)

#### Exercise 2c

Note that if B' obeys  $a = \bigcap_{b' \in B'} \operatorname{stab}(b') \in \Gamma$  and it is maximal under this property, then  $B' = \{b \in B | a \subseteq \operatorname{stab}(b)\}$ . For let  $b \in B$  such that  $a \subseteq \operatorname{stab}(b)$ . Then  $\bigcap_{b' \in B'} \operatorname{stab}(b') \cap \operatorname{stab}(b) \in \Gamma(\operatorname{because} \Gamma \text{ is a filter of subgroups})$ , so by maximality of B' we know  $b \in B'$ . So  $B' \subseteq \{b \in B | a \subseteq \operatorname{stab}(b)\}$ . Now note that for any  $b \in B'$  that  $a = \bigcap_{b' \in B'} \operatorname{stab}(b') \subseteq \operatorname{stab}(b)$ , so  $B' = \{b \in B | a \subseteq \operatorname{stab}(b)\}$ .

Now we find that for any  $x, y \in B'$  that  $\operatorname{stab}(x \wedge y) \supseteq \operatorname{stab}(x) \cap \operatorname{stab}(y) \supseteq a(\operatorname{as} \operatorname{if} g \in \operatorname{stab}(x) \cap \operatorname{stab}(y)$ , then  $g(x \wedge y) = gx \wedge gy = x \wedge y$ , so  $x \wedge y \in B'$ . Similarly  $x \vee y, x \Rightarrow y, x^* \in B'$ , so B' is a Boolean algebra (as these operations obey the required properties, as they do in B). (note that B' is a complete Boolean algebra by for  $X \subseteq B'$ ,  $\operatorname{stab}(\bigvee X) \supseteq \bigcap_{x \in X} \operatorname{stab}(x) \supseteq a$  so  $\bigvee X \in B'$ , and similarly  $\bigwedge X \in B'$ ). (1 point for proving that B' is a Boolean algebra)

Now we can show for any  $u \in V^{(B')}$  that  $\operatorname{stab}(u) \supseteq a$  by induction on  $V_{\alpha}^{(B')}$ . As let  $\alpha$  be an ordinal, and for any  $\beta < \alpha$  we know that for any  $u \in V_{\beta}^{(B')}$  that  $\operatorname{stab}(u) \supseteq a$ . Now let  $u \in V_{\alpha}^{(B')}$ . Then  $\operatorname{Fun}(u) \wedge \operatorname{ran}(u) \subseteq B' \wedge \exists \beta < \alpha \operatorname{dom}(u) \subseteq V_{\beta}^{(B')}$ . So let  $g \in a$ , then  $gu = \{\langle gx, g(u(x)) \rangle | x \in \operatorname{dom}(u) \}$ . By the induction hypothesis for any  $x \in \operatorname{dom}(u)$  we know that gx = x, and by definition of B' we know that for any  $b \in B'$  gb = b, so g(u(x)) = u(x). So  $gu = \{\langle x, u(x) \rangle | x \in \operatorname{dom}(u) \} = u$ . So  $\operatorname{stab}(u) \supseteq a$ .

Now by induction on ordinals  $\alpha$  we can find that  $V_{\alpha}^{(B')} \subseteq V_{\alpha}^{(\Gamma)}$ . As let  $\alpha$  be an ordinal, and for any  $\beta < \alpha$  we know that  $V_{\beta}^{(B')} \subseteq V_{\beta}^{(\Gamma)}$ . Now if  $x \in V_{\alpha}^{(B')}$  then  $\operatorname{Fun}(x) \wedge \operatorname{ran}(x) \subseteq B' \wedge \exists \beta < \alpha \operatorname{dom}(x) \subseteq V_{\beta}^{(B')}$ . Then by the induction hypothesis(and  $B' \subseteq B$ )  $\operatorname{Fun}(x) \wedge \operatorname{ran}(x) \subseteq B \wedge \exists \beta < \alpha \operatorname{dom}(x) \subseteq V_{\beta}^{(\Gamma)}$ , and by the previous part  $\operatorname{stab}(x) \supseteq a$  so  $\operatorname{stab}(x) \in \Gamma$ , so  $x \in V_{\alpha}^{(\Gamma)}$ . So by induction for any ordinal  $\alpha$  we know that  $V_{\alpha}^{(B')} \subseteq V_{\alpha}^{(\Gamma)}$ , so  $V^{(B')} \subseteq V^{(\Gamma)}$ . (1 point for correctly using induction)

# Exercise 3

#### Exercise 3a

Let  $g \in \operatorname{stab}(u)$ . Then

$$dom(gv) = \bigcup \{gdom(y)|y \in dom(u)\}$$
  
= 
$$\bigcup \{dom(gy)|y \in dom(u)\}$$
  
= 
$$\bigcup \{dom(y)|g^{-1}y \in dom(u)\}$$
  
= 
$$\bigcup \{dom(y)|y \in dom(gu)\} = dom(v)$$

using the property that  $\operatorname{dom}(gy) = \{gx | x \in \operatorname{dom}(y)\}$  for any  $y \in V^{(\Gamma)}$ .(.5 points) Now for  $x \in \operatorname{dom}(gv)$  we know that:

$$\begin{split} (gv)(x) &= g(v(g^{-1}(x))) \\ &= g[\![\exists y \in u[g^{-1}x \in y]]\!]^{\Gamma} \\ &= [\![\exists y \in gu[gg^{-1}x \in y]]\!]^{\Gamma} \\ &= [\![\exists y \in u[x \in y]]\!]^{\Gamma} = v(x) \end{split}$$

So  $g \in \operatorname{stab}(v)$ , so this completes the proof. (.5 points)

## Exercise 3b

Let  $g \in \operatorname{stab}(u)$ . Then if  $x \in \operatorname{dom}(gv)$ , then  $gx \in B^{\operatorname{dom}(u)} \cap V^{(\Gamma)}$ , so in other words  $gx \in V^{(\Gamma)}$  and  $\operatorname{dom}(gx) = \operatorname{dom}(u)$ . But if  $gx \in V^{(\Gamma)}$ , then  $x = g^{-1}gx \in V^{(\Gamma)}$ , and  $\operatorname{dom}(g^{-1}gx) = \operatorname{dom}(g^{-1}u) = \operatorname{dom}(u)$ . As the definition of v in this case is identical, we can just follow the proof(1 point).(if instead  $v(x) = [x \subseteq u$  which is what was required for the proof of thm 3.19, then we just write out definitions as in 3a)